

Formalizing Interoperability Testing

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Introduction

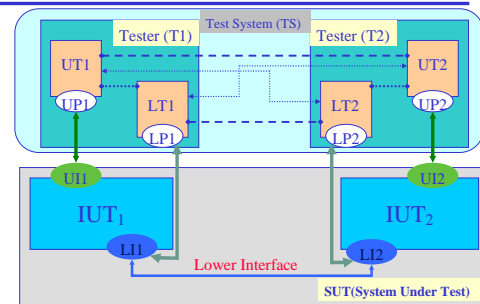
- **Interoperability test**: verification of the communication between the IUTs and of the provided services
- Some existing attempts to formalize interoperability but no precise characterization.
- **Interoperability test generation** : no method based on formal definitions
- Work presented here: **one-to-one context**
- **Future work: N>2**

Plan

- Introduction
- Interoperability testing architecture(s)
- Model for the formalization
- Interoperability definitions in the one-to-one context
- Interoperability test generation in the one-to-one context
- Conclusion and perspectives

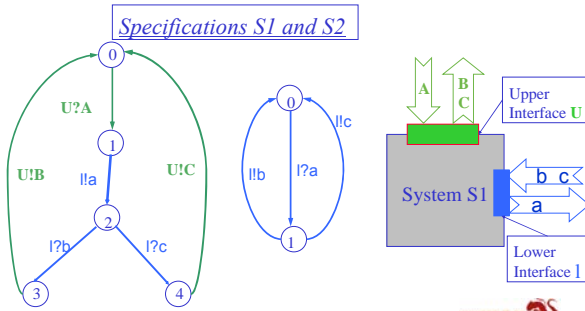
Test architectures

General interoperability testing architecture



Model for the formalization

Model of IOLTS and notations

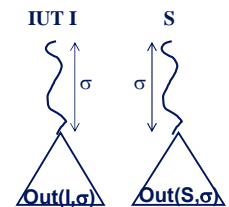


Model for the formalization

Some words about conformance testing

Conformance relation *ioconf*:
 $I \text{ ioconf } S \Leftrightarrow \forall \sigma \in \text{Traces}(S) \Rightarrow \text{Out}(I, \sigma) \subseteq \text{Out}(S, \sigma)$

Traces(S) : succession of events of the specification S
Out(S, σ) : set of possible outputs (sent messages) after the succession of events σ

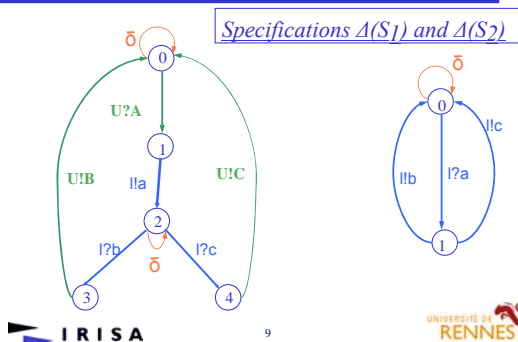


- Definition of *interoperability criteria* based on *ioconf*. Examples:
 - > "Global total Interoperability Criterion"

$$R_9(I_1, I_2) = \forall \sigma \in \text{Traces}(S_1 \parallel_{\mathcal{A}} S_2) \Rightarrow \text{Out}(I_1 \parallel_{\mathcal{A}} I_2, \sigma) \subseteq \text{Out}(S_1 \parallel_{\mathcal{A}} S_2, \sigma)$$
 R_9 equivalent to $I_1 \parallel_{\mathcal{A}} I_2$ *ioconf* $S_1 \parallel_{\mathcal{A}} S_2$
 - > "Unilateral total Interoperability Criterion"

$$R_7(I_1, I_2) = \forall \sigma_1 \in \text{Traces}(S_1), \forall \sigma \in \text{Traces}(I_1 \parallel_{\mathcal{A}} I_2) \sigma / \Sigma^1 = \sigma_1 \Rightarrow \text{Out}((I_1 \parallel_{\mathcal{A}} I_2) / \Sigma^1, \sigma) \subseteq \text{Out}(S_1, \sigma_1).$$
 R_7 equivalent to I_1 *ioconf* S_1 during the interaction of I_1 with I_2
- Different criteria considering the different testing architectures.

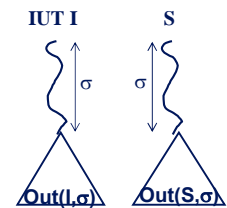
- Cases of non-interoperability not detected
 - Incorrect verdicts during tests based on these relations due to the absence of *quiescence management*.
- ⇒ **Quiescence**: deadlock, outputlock or livelock
- > Quiescence can be allowed in the specification(s)
 - > Represented with δ
 - > Considered as an *observable output event*



- Conformance relation *ioconf*

$$I \text{ ioconf } S = \forall \sigma \in \text{Traces}(S) \Rightarrow \text{Out}(I, \sigma) \subseteq \text{Out}(S, \sigma)$$
- Conformance relation *ioco* (with quiescence management)

$$I \text{ ioco } S = \forall \sigma \in \text{Traces}(\Delta(S)) \Rightarrow \text{Out}(I, \sigma) \subseteq \text{Out}(\Delta(S), \sigma)$$

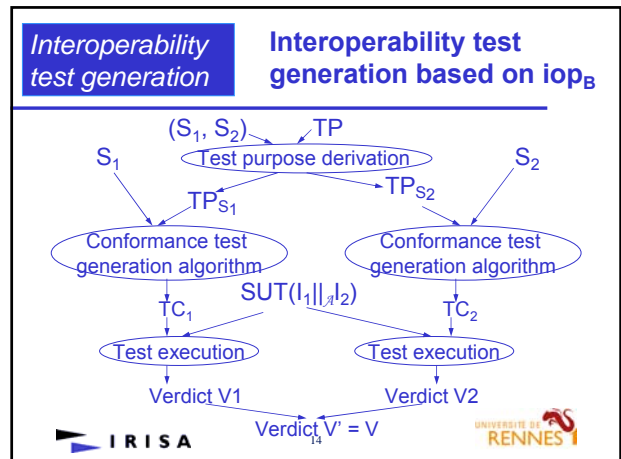
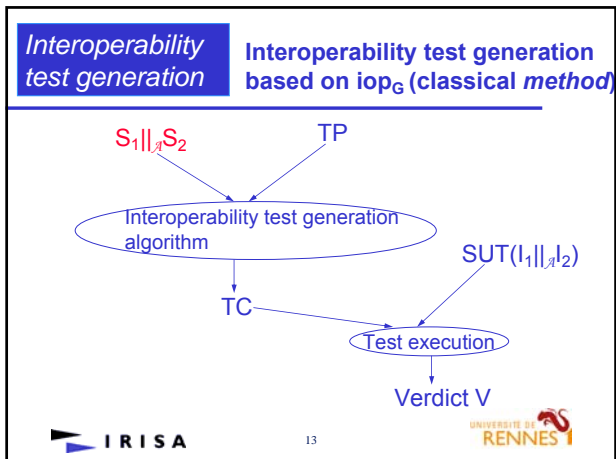


- *Global interoperability criterion*:

$$\text{iop}_G(I_1, I_2) = \forall \sigma \in \text{Traces}(S_1 \parallel_{\mathcal{A}} S_2) \Rightarrow \text{Out}(I_1 \parallel_{\mathcal{A}} I_2, \sigma) \subseteq \text{Out}(S_1 \parallel_{\mathcal{A}} S_2, \sigma)$$
- *Bilateral interoperability criterion*: $\text{iop}_B(I_1, I_2) = \forall \sigma_1 \in \text{Traces}(\Delta(S_1)), \sigma \in \text{Traces}(I_1 \parallel_{\mathcal{A}} I_2), \sigma / \Sigma^1 = \sigma_1 \Rightarrow \text{Out}((I_1 \parallel_{\mathcal{A}} I_2) / \Sigma^1, \sigma) \subseteq \text{Out}(\Delta(S_1), \sigma_1)$ and

$$\forall \sigma_2 \in \text{Traces}(\Delta(S_2)), \sigma \in \text{Traces}(I_1 \parallel_{\mathcal{A}} I_2), \sigma / \Sigma^1 = \sigma_2 \Rightarrow \text{Out}((I_1 \parallel_{\mathcal{A}} I_2) / \Sigma^1, \sigma) \subseteq \text{Out}(\Delta(S_2), \sigma_2).$$

- Non-interoperability due to a non-allowed output still detected
- More non-interoperability cases detected :
 - > One due to non-allowed quiescence in one of the IUT
 - > One due to incompatibility between output of one IUT and input of the other



- Interoperability test generation** Input-output causal dependency based method
- Previous methods only verify that *outputs (and quiescence)* are foreseen in the specifications
 - **Problem:** tester cannot verify if corresponding *inputs* are really executed (by the other IUT)
 - **Solution:** use of *causal dependencies* between an input and the possible output to conclude
 - **Contributions** of the method compared with the *iop-criteria-based* method:
 - > Interoperability test purpose expressiveness
 - > Pass and Inc verdict refinement
 - > Interoperability notion more taken into account
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- Conclusion and perspectives** Conclusion on one-to-one interoperability context
- Formal definitions (including quiescence management)
 - Interoperability test generation algorithms:
 - > Based on the formal definitions
 - > Input-output causal dependency based
 - Future work:
 - Implementation of these algorithms and test on real protocols
 - Generalization to $N>2$ implementations
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- Conclusion and perspectives** Interoperability with $N>2$ IUT
- $N>2$ specifications (we can have $S_i=S_j \forall (i,j)$ with $i \neq j$)
 - Interconnexion topologies:
 - Depends on the number of IUT
 - Some topologies are equivalent
 - Test architecture: depends on the topology
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